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Continuation-passing style

But first...



problem, boole?

Loose comparisons with ==

	TRUE	FALSE	1	0	-1	"1"	"0"	"-1"	NULL	array()	"php"	""
TRUE	TRUE	FALSE	TRUE	FALSE	TRUE	TRUE	FALSE	TRUE	FALSE	FALSE	TRUE	FALSE
FALSE	FALSE	TRUE	FALSE	TRUE	FALSE	FALSE	TRUE	FALSE	TRUE	TRUE	FALSE	TRUE
1	TRUE	FALSE	TRUE	FALSE	FALSE	TRUE	FALSE	FALSE	FALSE	FALSE	FALSE	FALSE
0	FALSE	TRUE	FALSE	TRUE	FALSE	FALSE	TRUE	FALSE	TRUE	FALSE	TRUE	TRUE
-1	TRUE	FALSE	FALSE	FALSE	TRUE	FALSE	FALSE	TRUE	FALSE	FALSE	FALSE	FALSE
"1"	TRUE	FALSE	TRUE	FALSE	FALSE	TRUE	FALSE	FALSE	FALSE	FALSE	FALSE	FALSE
"0"	FALSE	TRUE	FALSE	TRUE	FALSE	FALSE	TRUE	FALSE	FALSE	FALSE	FALSE	FALSE
"-1"	TRUE	FALSE	FALSE	FALSE	TRUE	FALSE	FALSE	TRUE	FALSE	FALSE	FALSE	FALSE
NULL	FALSE	TRUE	FALSE	TRUE	FALSE	FALSE	FALSE	FALSE	TRUE	TRUE	FALSE	TRUE
array()	FALSE	TRUE	FALSE	FALSE	FALSE	FALSE	FALSE	FALSE	TRUE	TRUE	FALSE	FALSE
"php"	TRUE	FALSE	FALSE	TRUE	FALSE	FALSE	FALSE	FALSE	FALSE	FALSE	TRUE	FALSE
""	FALSE	TRUE	FALSE	TRUE	FALSE	FALSE	FALSE	FALSE	TRUE	FALSE	FALSE	TRUE

'9223372036854775807' == '9223372036854775808'

Final project

$\langle \text{prog} \rangle ::= \langle \text{defs} \rangle \langle \text{exp} \rangle$

$\langle \text{defs} \rangle ::= \langle \text{def} \rangle \langle \text{defs} \rangle$
|

$\langle \text{def} \rangle ::= (\text{define } (\langle \text{var} \rangle \langle \text{formals} \rangle) \langle \text{exp} \rangle)$

$\langle \text{formals} \rangle ::= \langle \text{var} \rangle \langle \text{formals} \rangle$
|

$\langle \text{exp} \rangle ::= \langle \text{integer} \rangle$

$\langle \text{exp} \rangle ::= \langle \text{integer} \rangle$
 $\quad \quad \quad | \quad \langle \text{var} \rangle$

$\langle \text{exp} \rangle ::= \langle \text{integer} \rangle$
 $\quad | \langle \text{var} \rangle$
 $\quad | \#t \mid \#f$

$\langle \text{exp} \rangle ::= \langle \text{integer} \rangle$
 $\quad | \langle \text{var} \rangle$
 $\quad | \#t \mid \#f$
 $\quad | (\text{lambda } (\langle \text{formals} \rangle) \langle \text{exp} \rangle)$

```
<exp> ::= <integer>
        | <var>
        | #t | #f
        | (lambda (<formals>) <exp>)
        | (let ([<var> <exp>] ...) <exp>)
```

```
<exp> ::= <integer>
        | <var>
        | #t | #f
        | (lambda (<formals>) <exp>)
        | (let ([<var> <exp>] ...) <exp>)
        | (if <exp> <exp> <exp>)
```

```
<exp> ::= <integer>
        | <var>
        | #t | #f
        | (lambda (<formals>) <exp>)
        | (let ([<var> <exp>] ...) <exp>)
        | (if <exp> <exp> <exp>)
        | (set! <var> <exp>)
```

```
<exp> ::= <integer>
| <var>
| #t | #f
| (lambda (<formals>) <exp>)
| (let ([<var> <exp>] ...) <exp>)
| (if <exp> <exp> <exp>)
| (set! <var> <exp>)
| (begin <exp> ...)
```

```
<exp> ::= <integer>
| <var>
| #t | #f
| (lambda (<formals>) <exp>)
| (let ([<var> <exp>] ...) <exp>)
| (if <exp> <exp> <exp>)
| (set! <var> <exp>)
| (begin <exp> ...)
| (<exp> <exp> ...)
```

`<exp> ::= <integer>`
`| <var>`
`| #t | #f`
`| (lambda (<formals>) <exp>)`
`| (let ([<var> <exp>] ...) <exp>)`
`| (if <exp> <exp> <exp>)`
`| (set! <var> <exp>)`
`| (begin <exp> ...)`
`| (<exp> <exp> ...)`
`| (<prim> <exp> ...)`

`<exp> ::= <integer>`
`| <var>`
`| #t | #f`
`| (lambda (<formals>) <exp>)`
`| (let ([<var> <exp>] ...) <exp>)`
`| (if <exp> <exp> <exp>)`
`| (set! <var> <exp>)`
`| (begin <exp> ...)`
`| (<exp> <exp> ...)`
`| (<prim> <exp> ...)`
`| (call/cc <exp>)`

`<prim> ::= + | - | * | =
 | void`

(eval ...)

~~(every...)~~

Due: April 26 AoE

Next week

- Tuesday lecture: Online
- Thursday lecture: Flux!

Continuation-Passing Style

CPS

What is CPS?

***A style* of programming**

An intermediate form

Why CPS?

Web programming

Simplifies interpreters

Eliminates call/cc

Two rules

Functions never return

Arguments are atomic

$(f (g x))$

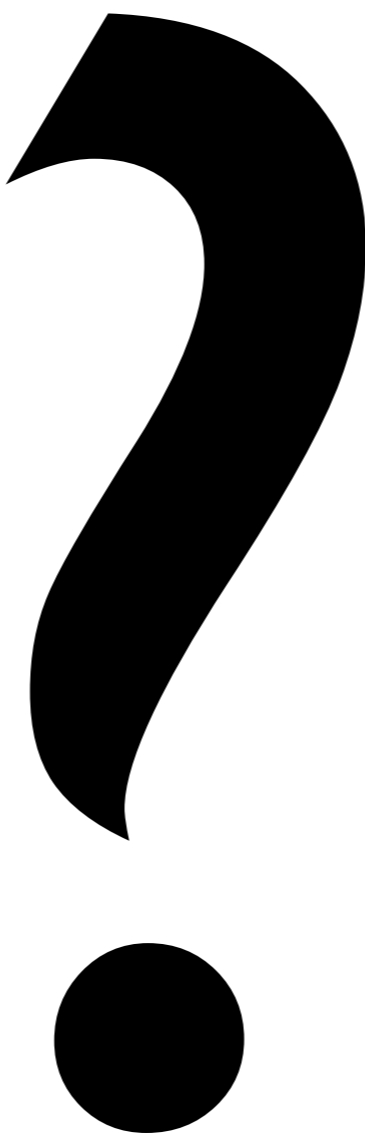
$(f(g))$



(lambda (x) x)

(lambda (x))





Pass callbacks

Pass continuations

$(\lambda (x) x)$

$(\lambda (x cc) (cc x))$

$(\lambda (x \text{ return})$
 $(\text{return } x))$

Example

```
(define (f n)
  (if (= n 0)
      1
      (* n (f (- n 1)))))
```

Example

```
(define (f n return)
  (if (= n 0)
      1
      (* n (f (- n 1)))))
```

Example

```
(define (f n return)
  (if (= n 0)
      (return 1)
      (* n (f (- n 1)))))
```

Example

```
(define (f n return)
  (if (= n 0)
      (return 1)
      (f (- n 1) (lambda (m)
                    (return (* n m))))))
```

Example

```
(define (f n return)
  (= $ n 0 (lambda (zero?)
    (if zero?
      (return 1)
      (- $ n 1 (lambda (sub)
        (f sub (lambda (mul)
          (* $ n mul return))))))))))
```

Example

```
(define (f a n return)
  (= $ n 0 (lambda (zero?)
    (if zero?
      (return a)
      (* $ a n (lambda (an)
        (- $ n 1 (lambda (n1)
          (f an n1 return))))))))))
```

Example

```
(define (fib n return)
  (<=$ n 0 (lambda (zero?)
    (if zero?
      (return n)
      (-$ n 1 (lambda (n1)
        (-$ n 2 (lambda (n2)
          (fib n1 (lambda (f1)
            (fib n2 (lambda (f2)
              (+$ f1 f2 return))))))))))))))
```

Naive transform

$expr ::= (\lambda (var) expr)$
| var
| $(expr expr)$

$$aexp ::= (\lambda (var_1 \dots var_n) cexp)$$

| var

$$cexp ::= (aexp_0 \dots aexp_n)$$

```
(define (T expr cont))
```

(define (T expr cont))

A term that invokes `cont` on the result of `expr`.

```
(define (T expr cont) `( ,cont ,expr))
```

```
(define (T expr cont)
  (match expr
```

```
(define (T expr cont)
  (match expr
    [(λ . _) `(,cont , expr )]
```

```
(define (T expr cont)
  (match expr
    [`(λ . ,_)      `(,cont ,(M expr))])
```

```
(define (M expr))
```

`(define (M expr))`

An equivalent expression that obeys rules of CPS.

```
(define (M expr)
  (match expr
```

```
(define (M expr)
  (match expr
    [ `(λ (, var) , expr)
```

```
(define (M expr)
  (match expr
    [ `(λ (, var) , expr)
      ; =>
      (define $k (gensym '$k))
```

```
(define (M expr)
  (match expr
    [ `(λ (, var) , expr)
      ; =>
      (define $k (gensym '$k))
      `(λ (, var , $k) ,(T expr $k))]
  )
)
```

```
(define (M expr)
  (match expr
    [`(λ (,var) ,expr)
     ; =>
     (define $k (gensym '$k))
     `(λ (,var , $k) ,(T expr $k))]
    [(? symbol?) #;=> expr]))
```

```
(define (T expr cont)
  (match expr
    [ `(λ . ,_)      `( ,cont , (M expr) ) ]
```

```
(define (T expr cont)
  (match expr
    [ `(λ . ,_)      `( ,cont , (M expr) ) ]
    [ (? symbol?)    `( ,cont , (M expr) ) ]
```

```
(define (T expr cont)
  (match expr
    [`(λ . ,_)      `( ,cont , (M expr))]
    [(? symbol?)    `( ,cont , (M expr))]
    [`( ,f ,e)
```

```
(define (T expr cont)
  (match expr
    [ `(λ . ,_) ` (,cont ,(M expr))]
    [ (? symbol?) ` (,cont ,(M expr))]
    [ `(,f ,e)
      ; =>
      (define $f (gensym '$f))
      (define $e (gensym '$e))
```

```
(define (T expr cont)
  (match expr
    [ `(λ . ,_) ` (,cont ,(M expr))]
    [ (? symbol?) ` (,cont ,(M expr))]
    [ `(,f ,e)
      ; =>
      (define $f (gensym '$f))
      (define $e (gensym '$e))
      (T f `(λ (,$f)
```

```

(define (T expr cont)
  (match expr
    [`(λ . ,_) ` (,cont ,(M expr))]
    [(? symbol?) ` (,cont ,(M expr))]
    [`(,f ,e)
     ; =>
     (define $f (gensym '$f))
     (define $e (gensym '$e))
     (T f `(λ (,$f)
              ,(T e `(λ (,$e)

```

```

(define (T expr cont)
  (match expr
    [`(λ . ,_) ` (,cont ,(M expr))]
    [(? symbol?) ` (,cont ,(M expr))]
    [`(,f ,e)
     ; =>
     (define $f (gensym '$f))
     (define $e (gensym '$e))
     (T f `(λ (,$f)
              ,(T e `(λ (,$e)
                       (,$f ,,$e ,cont))))))]
  ))

```

$(M \ ' \ (\lambda \ (x) \ x))$

$(\lambda (x \text{ \$k9}) (\text{\$k9 } x))$

(T ' (g a) 'halt)

$((\lambda (f9596)$

$((\lambda (e9597)$

$(f9596 e9597 halt)) a)) g)$

(g a halt)

call/cc?

(λ (f cc)

(f (λ (x cc*)

(cc x))

cc)))

Net effect?

CESK

CES

Braintaser

(call/cc call/cc)

See blog for more!